

Convergence of Iterative Hard Thresholding Variants with Application to Asynchronous Parallel Methods for Sparse Recovery

Jamie Haddock

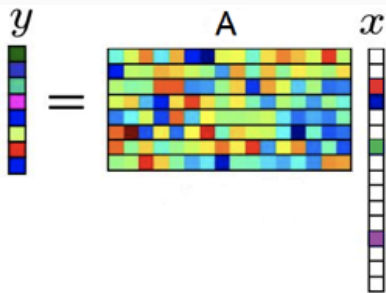
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UCLA



joint with Deanna Needell, Nazanin Rahnavard, and Alireza Zaezadeh

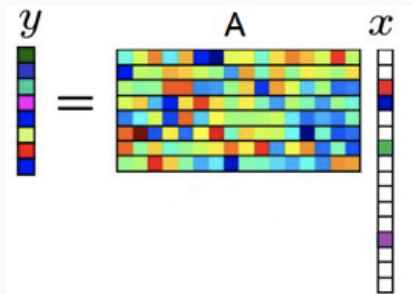
Sparse Recovery Problem



Sparse Recovery: reconstruct approximately sparse $\mathbf{x} \in \mathbb{R}^N$ from few nonadaptive, linear, and noisy measurements, $\mathbf{y} = \mathbf{A}\mathbf{x} + \mathbf{e}$

$\mathbf{A} \in \mathbb{R}^{m \times N}$: measurement matrix
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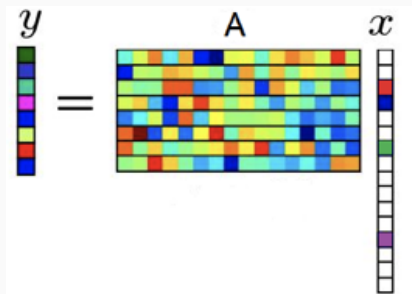
Approach:

$$\min_{\mathbf{x} \in \mathbb{R}^N} \|\mathbf{x}\|_1 \text{ s.t. } \|\mathbf{A}\mathbf{x} - \mathbf{y}\| \leq \epsilon$$

or

$$\min_{\mathbf{x} \in \mathbb{R}^N} \frac{1}{m} \|\mathbf{A}\mathbf{x} - \mathbf{y}\|^2 \text{ s.t. } \|\mathbf{x}\|_0 \leq s$$

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Applications:

- ▷ image reconstruction
- ▷ hyper spectral imaging
- ▷ wireless communications
- ▷ analog to digital conversion

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$$\text{IHT: } \mathbf{x}^{(n+1)} = H_k(\mathbf{x}^{(n)} + \mathbf{A}^T(\mathbf{y} - \mathbf{A}\mathbf{x}^{(n)}))$$

Algorithm 1 StoIHT Algorithm [22]**input:** $s, \gamma, p(i)$, and stopping criterion**initialize:** x^1 and $t = 1$ **repeat****randomize:** select $i_t \in [M]$ with probability $p(i_t)$ **proxy:** $b^t = x^t + \frac{\gamma}{Mp(i_t)} A_{b_{i_t}}^* (y_{b_{i_t}} - A_{b_{i_t}} x^t)$ **identify:** $\Gamma^t = \text{supp}_s(b^t)$ **estimate:** $x^{t+1} = b_{\Gamma^t}^t$
 $t = t + 1$ **until** halting criterion *true***output:** $\hat{x} = x^t$

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¹Nguyen, Needell, Woolf, IEEE Transactions on Information Theory '17

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Challenge: objective of $\min_{\mathbf{x} \in \mathbb{R}^N} \frac{1}{m} \|\mathbf{Ax} - \mathbf{y}\|^2$ s.t. $\|\mathbf{x}\|_0 \leq s$ is dense in \mathbf{x}

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- ▷ likely that same non-zero entries are updated from one iteration to the next
- ▷ a slow core could easily “undo” the progress of previous updates by faster cores

Algorithm 2 Asynchronous StoIHT Iteration

Each core performs the following at each iteration. The tally vector ϕ is available to each core.

randomize: select $i_t \in [M]$ with probability $p(i_t)$

proxy: $b^t = x^t + \frac{\gamma}{Mp(i_t)} A_{b_{i_t}}^* (y_{b_{i_t}} - A_{b_{i_t}} x^t)$

identify: $\Gamma^t = \text{supp}_s(b^t)$

$$\tilde{T}^t = \text{supp}_s(\phi)$$

estimate: $x^{t+1} = b_{\Gamma^t \cup \tilde{T}^t}^t$

update tally: $\phi_{\Gamma^t} = \phi_{\Gamma^t} + t$

$$\phi_{\Gamma^{t-1}} = \phi_{\Gamma^{t-1}} - (t - 1)$$

$$t = t + 1$$

²Needell, Woolf, Proc. Information Theory and Applications '17

Bayesian Asynchronous StoIHT

Require: Number of subproblems, M , and probability of selection $p(B)$.

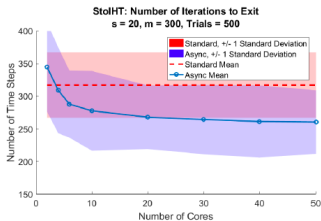
The reliability score distribution parameters, $\hat{\beta}_i^1$ and $\hat{\beta}_i^0$, and the tally scores parameters, \hat{a}_n^1 and \hat{a}_n^0 , are available to each processor.

Each processor performs the following at each iteration:

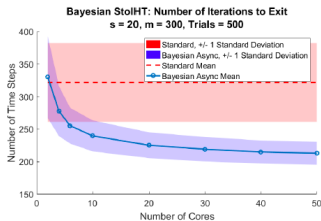
- 1: **randomize:** select $B_t \in [M]$ with probability $p(B_t)$
- 2: **proxy:** $\mathbf{b}^{(t)} = \mathbf{x}^{(t)} + \frac{\gamma}{Mp(B_t)} \mathbf{A}_{B_t}^* (\mathbf{y}_{B_t} - \mathbf{A}_{B_t} \mathbf{x}^{(t)})$
- 3: **identify:** $\hat{S}^{(t)} = \text{supp}_s(\mathbf{b}^{(t)})$ and $\tilde{T}^{(t)} = \text{supp}_s(\phi)$
- 4: **estimate:** $\mathbf{x}^{(t+1)} = \mathbf{b}^{(t)}_{\hat{S}^{(t)} \cup \tilde{T}^{(t)}}$
- 5: **repeat**
- 6: update $\mathbb{E}_{\mathbb{Q}\{u_{ni}\}} \{u_{ni}\} = \mathbb{Q}\{u_{ni} = 1\}$
- 7: update $\hat{\beta}_i^1$ and $\hat{\beta}_i^0$, \hat{a}_n^1 and \hat{a}_n^0
- 8: **until** convergence
- 9: update ϕ
- 10: $t = t + 1$

²Zaemzadeh, H., Rahnavard, Needell, Proc. 49th Asilomar Conf. on Signals, Systems and Computers '18

Experimental Convergence



(a) StoIHT



(b) Bayesian StoIHT

Figure 2: Number of time steps until convergence versus number of cores used in (a) asynchronous StoIHT method and (b) Bayesian asynchronous StoIHT. Half of the cores are *slow* and complete an iteration only once out of every four time steps.

First step: analyze IHT variant running on each node of parallel system

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Non-Symmetric Isometry Property:

$$(1 - \beta_k)\|\mathbf{z}\|_2^2 \leq \|\mathbf{A}\mathbf{z}\|_2^2 \leq \|\mathbf{z}\|_2^2 \text{ for all } k\text{-sparse } \mathbf{z}$$

Theorem (H., Needell, Zaemzadeh, Rahnavard '19+)

If \mathbf{A} has the non-symmetric restricted isometry property with $\beta_{3k+2\tilde{k}} < \frac{1}{8}$, then in iteration n , the IHT_{k, \tilde{k}} algorithms with input observations $\mathbf{y} = \mathbf{Ax} + \mathbf{e}$ recover the approximation $\mathbf{x}^{(n)}$ with

$$\|\mathbf{x} - \mathbf{x}^{(n)}\| \leq 2^{-n} \|\mathbf{x}^k\| + 5\|\mathbf{x} - \mathbf{x}^k\| + \frac{4}{\sqrt{k}} \|\mathbf{x} - \mathbf{x}^k\|_1 + 4\|\mathbf{e}\|.$$

Theorem (H., Needell, Zaemzadeh, Rahnavard '19+)

Suppose the signal \mathbf{x} has constant values on its support, and the \tilde{k} indices selected (non-greedily) by the $IHT_{k,\tilde{k}}$ algorithm each lie uniformly in the support of \mathbf{x} with probability p . If \mathbf{A} has the non-symmetric restricted isometry property with $\beta_{3k+2\tilde{k}} < \frac{1}{8}$, then in iteration n , the $IHT_{k,\tilde{k}}$ algorithms with input observations $\mathbf{y} = \mathbf{A}\mathbf{x} + \mathbf{e}$ recover the approximation $\mathbf{x}^{(n)}$ with

$$\begin{aligned}\mathbb{E}_{\tilde{k}} \|\mathbf{x} - \mathbf{x}^{(n)}\| &\leq 2^{-n} \|\mathbf{x}\| + 5\mathbb{E}_{\tilde{k}} \|\mathbf{x} - \tilde{\mathbf{x}}^{(n)}\| \\ &\quad + \frac{4}{\sqrt{k}} \mathbb{E}_{\tilde{k}} \|\mathbf{x} - \tilde{\mathbf{x}}^{(n)}\|_1 + 4\|\mathbf{e}\| \\ &\leq 2^{-n} \|\mathbf{x}\| + \left(5\alpha + \frac{4\alpha}{\sqrt{k}}\right) \|\mathbf{x}\|_1 + 4\|\mathbf{e}\|\end{aligned}$$

where $\alpha = \left(\frac{|\text{supp}(\mathbf{x})| - k}{|\text{supp}(\mathbf{x})|}\right) \left(\frac{|\text{supp}(\mathbf{x})| - p\tilde{k}}{|\text{supp}(\mathbf{x})|}\right)$.

Experimental Convergence of IHT _{k, \tilde{k}}

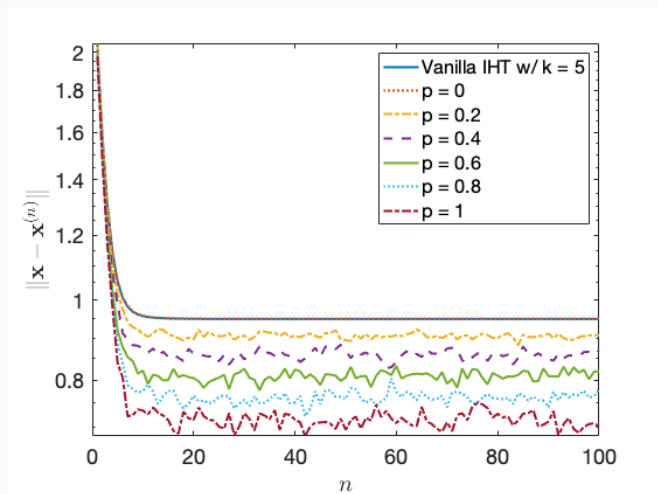
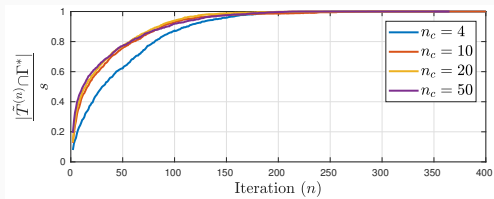
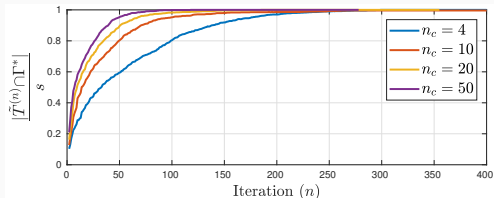


Figure 1: Plot of error $\|\mathbf{x} - \mathbf{x}^{(n)}\|$ vs. iteration for 100 iterations of IHT _{k, \tilde{k}} with various probabilities p that the \tilde{k} indices lie in $\text{supp}(\mathbf{x})$.

Rate of Support Intersection



(a)



(b)

Figure 2: The rate at which the shared indices between nodes lie in the true support of signal \mathbf{x} for iterations of (a) ASolHT and (b) BASolHT.

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- ▷ provided heuristic for why asynchronous versions of StoIHT converge faster than non-parallelized version

- ▷ analyze StoIHT _{k, \tilde{k}}
- ▷ extend to non-heuristic analysis of Asynchronous StoIHT

Questions?

- [1] J. Haddock, D. Needell, N. Rahnavard, and A. Zaeemzadeh. **Convergence of iterative hard thresholding variants with application to asynchronous parallel methods for sparse recovery.** In Proc. Asilomar Conf. Sig. Sys. Comp., 2019.
- [2] Deanna Needell and Tina Woolf. **An asynchronous parallel approach to sparse recovery.** In 2017 Information Theory and Applications Workshop (ITA), pages 1–5. IEEE, 2 2017.
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- [4] A. Zaeemzadeh, J. Haddock, N. Rahnavard, and D. Needell. **A Bayesian approach for asynchronous parallel sparse recovery.** In Proc. Asilomar Conf. Sig. Sys. Comp., 2018.